Algorithms

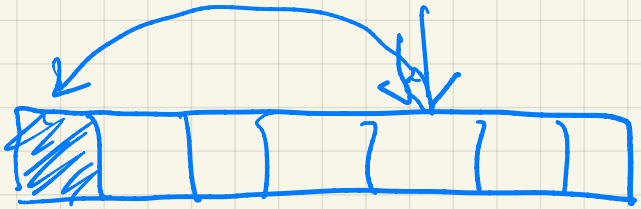
Book: Leiserson MIT

CLRS

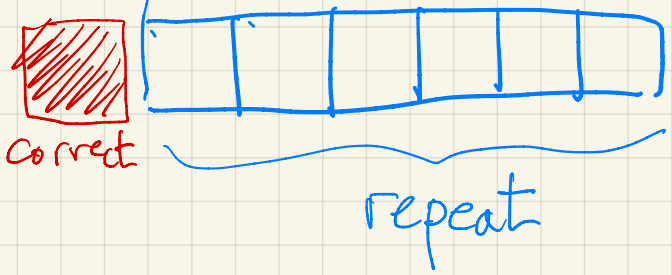
Chapter 2. (Sorting)

Sorting Algorithm:

- Bubble Sort
- Selection Sort

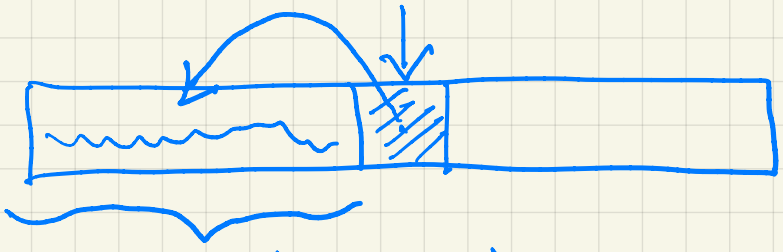


- Find smallest (scan entire array)
- switch with first

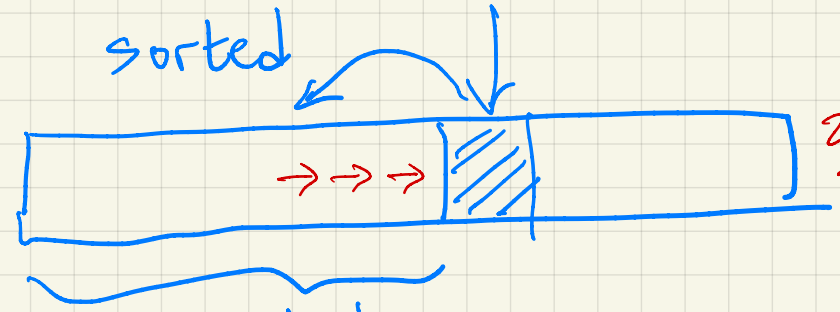


Amount of work: not sensitive to original order

Insertion Sort

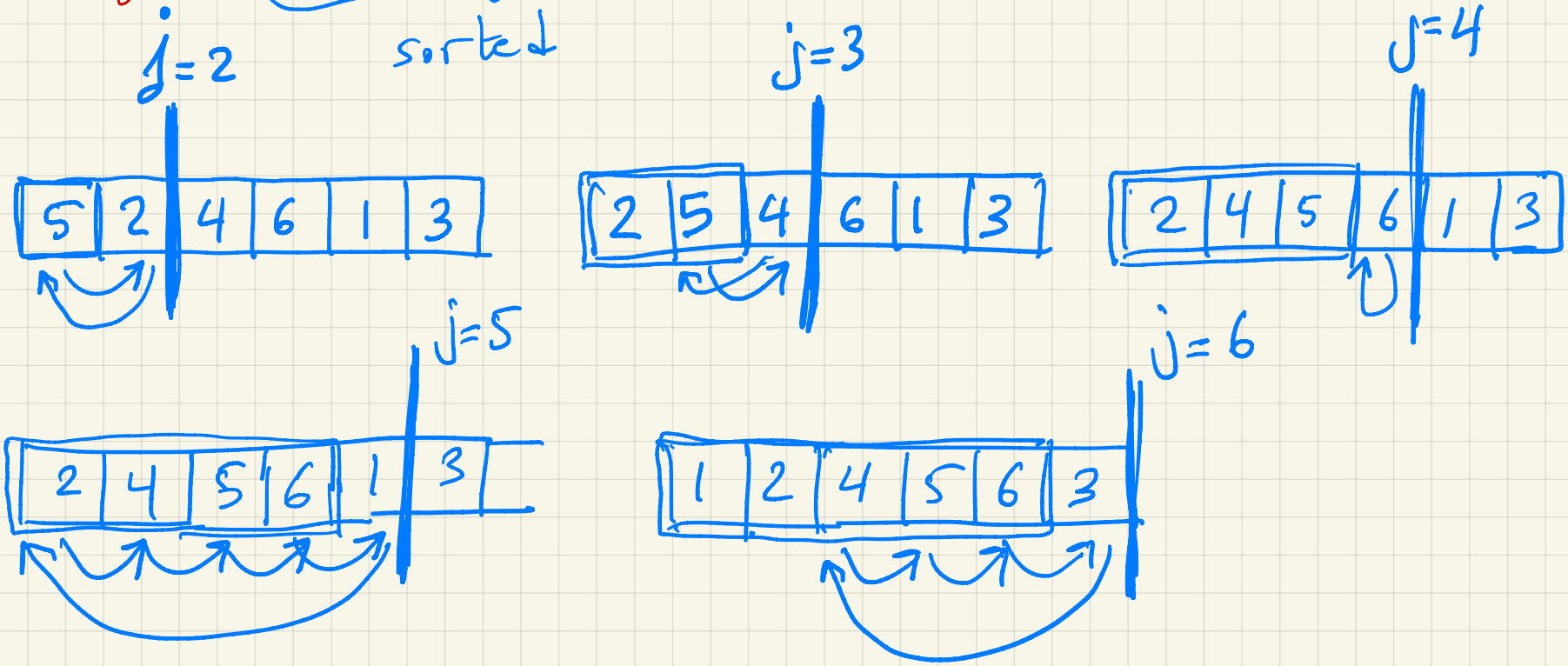


1) Assume $A[1 \cdots j-1]$ sorted



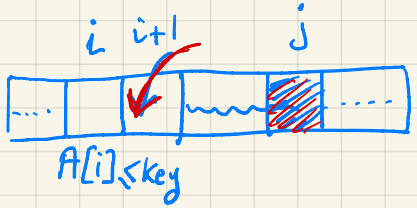
2) Insert $A[j]$ in the correct position

Example:



		Costs	# times
	Insetion Sort (A)		
→	for $j \leftarrow 2$ to n	1	n
	do key $\leftarrow A[j]$	1	$n-1$
	▷ Insert $A[j]$ into the sorted sequence $A[1 \dots j-1]$		
	$i \leftarrow j-1$	1	$n-1$
shifting process	while $i > 0$ and $A[i] > key$	1	$\sum_{j=2}^{n} t_j$
shifting process	do $A[i+1] \leftarrow A[i]$	1	$\sum_{j=2}^{n} (t_j - 1)$
shifting process	$i \leftarrow i-1$	1	$\sum_{j=2}^{n} (t_j - 1)$
insertion	$A[i+1] \leftarrow key$	1	$n-1$

After while loop



Loop invariant: sub array $A[1 \dots j-1]$ at the start of each iteration consists of all the elements originally in $A[1 \dots j-1]$ in sorted order

Initialization: $A[1]$ consists of all elements that are originally in $A[1]$ and sorted.

Maintenance: ✓ think about it

Termination: $j = n+1$

$A[1 \ldots j-1] = A[1 \ldots n]$ consists of all elements that were originally in $A[1 \ldots n]$ in sorted order.

Total time of insertion sort.

$$n + (n-1) + (n-1) + (n-1) + \sum_{j=2}^{n} t_j + 2\sum_{j=2}^{n}(t_j - 1)$$

$$n + n - 1 + 3\sum_{j=2}^{n} t_j \qquad 2\sum_{j=2}^{n}(-1) \qquad -2(n-1)$$

$$\boxed{2n - 1 + 3\sum_{j=2}^{n} t_j}$$