

Introduction to the Theory of Computation

Homework 1

Due 9/8/2017

These are practice exercises

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Problem 0: Readings

(a) Read Chapter 0 in Sipser's book.

(b) Read Chapter 1 in Sipser's book.

Problem 1

The formal definition of a DFA M is $(\{q_1, q_2, q_3, q_4, q_5\}, \{u, d\}, \delta, q_3, \{q_3\})$, where δ is given by the following table:

	u	d
q_1	q_1	q_2
q_2	q_1	q_3
q_3	q_2	q_4
q_4	q_3	q_5
q_5	q_4	q_5

Give the state diagram of this machine.

Problem 2 (do as many as you want)

For each of the following languages, construct a DFA that accepts the language.

Assume $\Sigma = \{0, 1\}$.

- $\{w \mid \text{the length of } w \text{ is divisible by three}\}$
- $\{w \mid w \text{ contains at least five } 1s\}$
- $\{w \mid w \text{ contains the substring } 1011\}$
- $\{w \mid w \text{ contains at least two } 1s \text{ and at most two } 0s\}$
- $\{w \mid w \text{ contains an odd number of } 1s \text{ or exactly two } 0s\}$

- $\{w|w \text{ begins with } 1 \text{ and ends with } 0\}$
- $\{w| \text{ every odd position in } w \text{ is } 1\}$
- $\{w|w \text{ has length at least } 3 \text{ and its third symbol is } 0\}$
- $\{\epsilon, 0\}$

Problem 3 (do as many as you want)

For each of the following languages, construct an NFA, with the specified number of states, that accepts the language. Again, $\Sigma = \{0, 1\}$.

- $\{w|w \text{ ends with } 10\}$ with 3 states
- $\{w|w \text{ contains the substring } 1011\}$ with five states
- $\{w|w \text{ contains an odd number of } 1s \text{ or exactly two } 0s\}$ with six states
- $\{w|w \text{ contains the substring } 11001\}$ with any number of states
- $\{w|w \text{ contains at least two } 1s \text{ and does not end with } 10\}$ with any number of states
- $\{w|w \text{ begins with } 1 \text{ or ends with } 0\}$ with any number of states

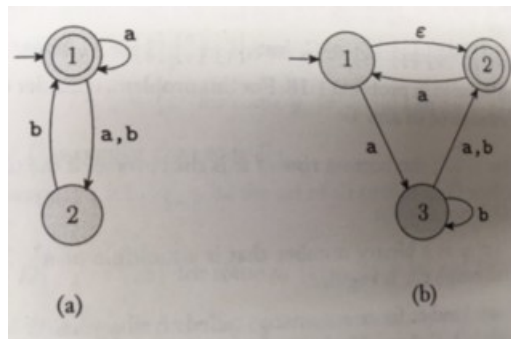
Problem 4

(a) Show that if M is a DFA that accepts language B , exchanging the accept and non-accept states in M yields a new DFA that accepts the complement of B , and conclude that the class of regular languages is closed under complement.

(b) Show, by giving a counter example, that if M is an NFA that accepts language C , exchanging the accept and non-accept states of M does not necessarily yield a new NFA that accepts the complement of C . Is the class of languages accepted by NFAs closed under complement? Explain your answer.

Problem 5

Convert the following NFAs to equivalent DFAs.



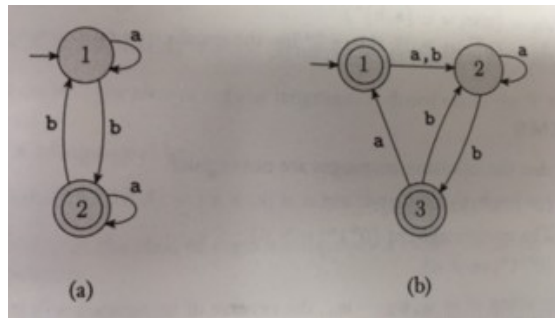
Problem 6

Convert the following regular expressions to NFAs.

- $(0 \cup 1)^*000(0 \cup 1)^*$
- $((00)^*(11) \cup 01)^*$
- \emptyset^*

Problem 7

Convert the following finite automata to regular expressions.

**Problem 8**

Use the pumping lemma to show that the following languages are not regular:

- $\{0^n1^n2^n | n \geq 0\}$
- $\{ww | w \in \{a, b\}^*\}$
- $\{1^{2^n} | n \geq 0\}$